Program: BE Computer Engineering Curriculum Scheme: Revised 2016 Examination: Final Year Semester VIII Course Code: CSC802 and Course Name: Distributed Computing

Time: 1 hour Max Marks:50

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Critical section Mutual Exclusion
Mutual Exclusion
Message passing
The is used to store a constant value that is decided based on the frequency of oscillation of the quartz of

	z crystal
	ant register
ALU	
	ich algorithm, One process is elected as the coordinator.
	buted mutual exclusion algorithm
	•
	alized mutual exclusion algorithm
	r Algorithm
Lamp	ort algorithm
If a ar	nd b are two events within the same process and a occurs before b, then clock of a (C(a)) is alwaysthen clock of
lesser	
greate	er
equal	
Cant	
	tributed systems, a logical clock is associated with .
	instruction
_	
	process
	register
each /	
	is a physical clock synchronization Algorithm
Cristia	an
vecto	r
Lamp	ort
Ring	
_	event a and b are parallel, then .
	b monotonic
C(a) =	
C(a) >	
C(a) <	
To en	force two functions are provided enter-critical and exit-critical, where each function takes as an argur
name	of the resource that is the subject of competition.
Deadl	ock
Starva	ation
	al exclusion
	ronization
_	is used to provide security to the shared resources, processes & channels used for their interactions.
The _	
_	action model
	model
	tectural model
Secur	ity model
The _	model is required in order to build systems with predictable behavior in case of fault
Intera	action model
Fault	model
	tectural model
	ity model
	defines the way in which the components of the system interact with each other & mapped onto an underlying networl
	, , , , , , , , , , , , , , , , , , , ,
	onent.
	rction model
	model
Archit	tectural model
Secur	ity model
	allows multiple computers with diverse hardware, operating system & network to solve complex problem
Cluste	er computing
Grid o	computing
	alized system
	er and gri computing
	er and gri computing In the following option which is not main focus of Distributed System
Availa	
Reliab	,
	lific Performance
	urce sharing
Secur	ity for information resource does not include
	ability
	urrency
Availa	dentiality
Availa Concu	
Availa Concu Confid	,
Availa Concu Confid Integr	ity
Availa Concu Confid Integr Which	rity n among below is not a technique commonly used for scaling in Distributed System
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What what is a second of
Workstation model
Workstation -server model
Processor-pool model Client-server model
The V-system in distributed computing system is based on .
Workstation model
Workstation -server model
Processor-pool model
Client-server model
The response time is short since process migration is not required in
Workstation model
Workstation -server model
Processor-pool model
Client-server model
refers to degree of tolerance against component failure and errors
Performance
Availability
Integrity
Reliability
Which among below is not Distributed system models are
Architectural model
Interaction model
Fault model
Performance model
is computing form where heterogeneity is key advantage.
Cluster computing
High performance computing
Grid computing
Cluster & grid computing
is software architecture used to build a distributed system from a network connected by high speed network.
DOS NOS
NOS Atidal pugas
Middleware On continuo system
Operating system The transparency that enables multiple instances of resources to be used, is called transparency
Concurrency
Performance
Scaling
Replication
implies that the user need not be aware whether the resource is in volatile memory or on the disk.
Relocation Transparency
Location Transparency
Migration Transparency
Persistence Transparency
Which is not a characteristic in Distributed Operating System
Enabling Interprocess communication
Uniform process management
Local control management
Different kernel implementation
A system enables a distributed system to behave like a virtual uniprocessor, even though the system operates on
collection of machines
DOS
NOS
Middleware
Operating system
The type of communication where a message is stored by the communication system only as long as the sending and receiving application
are executing.
Transient communication
Persistent communication Synchronous communication
Asynchronous communication
The type of communication where a message that has been submitted for transmission is stored by the communication system as long as
it takes to deliver it to the receiver
Transient communication
Persistent communication
Synchronous communication
Synchronous communication Asynchronous communication
·
The type of communication where a sender continues its execution immediately after it has submitted its message for transmission
·
The type of communication where a sender continues its execution immediately after it has submitted its message for transmission

	ommunication
•	communication
The type of c	ommunication where the sender is blocked until its message is stored in a local buffer at the receiving host,
delivered to th	e receiver.
Transient com	munication
Persistent con	
,	ommunication
,	communication
	is an example of
	nchronous communication
	chronous communication
	chronous communication
UDP is an exar	hronous communication
	nchronous communication
	chronous communication
	chronous communication
	hronous communication
	ating system on the server machine passes the incoming packets to the
Server stub	
Client stub	
Client operatir	g system
Client process	
	of protocols used in a particular system is called a protocol suite or
Protocol collec	tor
Protocol list	
Protocol stack	
Protocol queu	
-	a server makes it possible for a client to locate the server and
Bind to it	
Listen to it Refer to it	
Store to it	
	tion in distributed systems is based on sending and receiving (low level) messages because of
	f shared memory
	of heterogeneity
	f synchronization
	of complexities
<u> </u>	the following is not performed be RPC Runtime
Retransmissio	1
Routing	
Marshalling	
Encryption	
In RPC,	handles transmission of messages across the network between client and the server machine
Rpc transmiss	on
Rpc runtime	ation
Rpc communion Rpc interface	alui
•	omputer program to cause a subroutine to execute in
Its own addres	
Another addre	I .
	ddress space and another address space
Applications a	
• •	offers intermediate-term storage capacity for messages, without requiring either the sender or receiver to be ac
•	mission is suitable for
	nchronous communication
	chronous communication
Persistent syn	chronous communication
Transient sync	hronous communication
Method invoc	ations between objects in the same process are
Static method	invocations
	od invocations
Local method	
	hod invocations
	nat can be used throughout a distributed system to refer to a particular unique remote object is
Remote objec	
Remote objec	
Remote objec	,
Remote objec	
	object has a remote interface that specifies which of its
Mothade can	pe invoked remotely
	n be used remotely
Parameters ca	
Parameters ca Methods can l	n be invoked locally
Parameters ca Methods can l Parameters ca	
Parameters ca Methods can l Parameters ca An RPC (remo	te procedure call) is initiated by the
Parameters ca Methods can I Parameters ca An RPC (remo Server	te procedure call) is initiated by the
Parameters ca Methods can I Parameters ca An RPC (remo Server Client	
Parameters ca Methods can I Parameters ca An RPC (remo Server	

An A r	
Αr	reation
	action
An	eaction
	invitation
Int	er Process Communication takes place via
Sha	ared memory and message passing
De	centralized memory and message passing
Sha	ared memory and message replicating
De	centralized memory and message replicating
Wł	nich failures Response lies outside a specified time interval
-	ning Failure
_	nission failure
-	ash Failure
	oitrary Failure
, (i k	and y runare
clic	ent-centric consistency provides guarantees for a concerning the consistency of accesses to a data store by that cli
	,, , , , , , , , , , , , , , , , , , , ,
	gle Client
	ultiple Client
	gle Server
-	ultiple Server
Wł	nich of the following is not example of Information redundancy
pai	rity
che	ecksum
Ha	mming codes
Ro	llback if transaction aborts
In	all processes see only those memory reference operations in the correct order that are potentially causally rela
_	ntinuous Consistency
-	usal Consistency
-	quential Consistency
	entual Consistency
	nich of the following is not Message Ordering Multicast
_	
	ordered multicasts
	O-ordered multicasts
_	usally-ordered multicasts
LIF	O-ordered multicasts
Wł	nich multicast delivers messages so that potential causality between different messages is preserved
Un	ordered multicasts
FIF	O-ordered multicasts
Ca	usally-ordered multicasts
To	tally-ordered multicasts
Wł	nich of the following is not comes under Dependability
-	ailability
-	liability
-	,
	sintainability
-	dundancy
-	system goes down for one millisecond every hour, then how much availability it has?
	09 percent
_	9999 percent
00.	99 percent
99.	3 percent
_	used in replication of stocks
Nu	merical deviation
Sta	lleness deviation
Ord	dering deviation
	ne deviation
Tin	is often used for Web caches
Tin	
	merical deviation
Nu	merical deviation
Nu Sta	leness deviation
Nu Sta Ord	leness deviation dering deviation
Nu Sta Ord	leness deviation dering deviation ne deviation
Nu Sta Oro Tin Wh	lleness deviation dering deviation ne deviation nich of the following is disadvantage of active replication
Nu Sta Ord Tin Wh	lleness deviation dering deviation ne deviation nich of the following is disadvantage of active replication nsistent replicas
Nu Sta Ord Tin Wh Co	lleness deviation dering deviation ne deviation nich of the following is disadvantage of active replication nsistent replicas nimal bandwidth costs
Nu Sta Ord Tin Wh Co Mi	lleness deviation dering deviation ne deviation nich of the following is disadvantage of active replication nsistent replicas nimal bandwidth costs e of the parameters is small
Nu Sta Ord Tin Wh Co Mi Size	eleness deviation dering deviation ne deviation nich of the following is disadvantage of active replication nisistent replicas nimal bandwidth costs e of the parameters is small occasing power
Nu Sta Ord Tin Wh Co Mi Size	lleness deviation dering deviation ne deviation nich of the following is disadvantage of active replication nsistent replicas nimal bandwidth costs e of the parameters is small
Nu Sta Orc Tin Wh Co Mi Siz Pro	eleness deviation dering deviation ne deviation nich of the following is disadvantage of active replication nisistent replicas nimal bandwidth costs e of the parameters is small occasing power
Nu Sta Orc Tin Wh Col Mil Size Pro The	eleness deviation dering deviation ne deviation nich of the following is disadvantage of active replication nisistent replicas nimal bandwidth costs e of the parameters is small decessing power e effect of a write operation by a process on data item x will always be seen by a successive read operation on x by the same
Nu Sta Orc Tin Wh Co Min Siz Pro The pro Rea	eleness deviation dering deviation ne deviation nich of the following is disadvantage of active replication nsistent replicas nimal bandwidth costs e of the parameters is small pressing power e effect of a write operation by a process on data item x will always be seen by a successive read operation on x by the same pocess.
Nu Sta Oro Tin Wh Co Mi Size Pro The pro Rea Mo	eleness deviation dering deviation ne deviation nich of the following is disadvantage of active replication nisistent replicas nimal bandwidth costs e of the parameters is small decessing power e effect of a write operation by a process on data item x will always be seen by a successive read operation on x by the same coess. ad-your-writes contonic-writes
Nu Sta Orc Tin Wh Co Mi Size Pro The pro Rea Mo	eleness deviation dering deviation ne deviation nich of the following is disadvantage of active replication nisistent replicas nimal bandwidth costs e of the parameters is small nocessing power e effect of a write operation by a process on data item x will always be seen by a successive read operation on x by the same nocess. ad-your-writes notonic-writes notonic-Reads
Nu Sta Ord Tin Wh Co Mi Siz Pro The pro Rea Mc	elleness deviation dering deviation ne deviation nich of the following is disadvantage of active replication nisistent replicas nimal bandwidth costs e of the parameters is small ocessing power e effect of a write operation by a process on data item x will always be seen by a successive read operation on x by the same nocess. ad-your-writes nontonic-writes notonic-writes notonic-Reads ittes-follow-reads
Nu Sta Orc Tin Wh Co Mi Size Pro The pro Rea Mo Wr	eleness deviation dering deviation ne deviation nich of the following is disadvantage of active replication nisistent replicas nimal bandwidth costs e of the parameters is small decessing power e effect of a write operation by a process on data item x will always be seen by a successive read operation on x by the same objects. ad-your-writes conotonic-writes conotonic-Reads cites-follow-reads nich Consistency is used to guarantee that users of a network newsgroup see a posting of a reaction to an article only after the
Nu Sta Orc Tin Wh Co Mi Size Pro The pro Rea Mc Wr Wr hav	eleness deviation dering deviation ne deviation nich of the following is disadvantage of active replication nisistent replicas nimal bandwidth costs e of the parameters is small decessing power e effect of a write operation by a process on data item x will always be seen by a successive read operation on x by the same occess. ad-your-writes onotonic-writes onotonic-writes onotonic-Reads ites-follow-reads nich Consistency is used to guarantee that users of a network newsgroup see a posting of a reaction to an article only after the verseen the original article
Nu Sta Orc Tin Wh Siz Pro The pro Mo Wr Wh	eleness deviation dering deviation ne deviation nich of the following is disadvantage of active replication nisistent replicas nisistent replicas nimial bandwidth costs e of the parameters is small coessing power e effect of a write operation by a process on data item x will always be seen by a successive read operation on x by the same occss. ad-your-writes onotonic-writes onotonic-writes onotonic-writes intes-follow-reads ich Consistency is used to guarantee that users of a network newsgroup see a posting of a reaction to an article only after the verseen the original article ad-your-writes
Nu Sta Orc Tin Wh Co Mi Size Pro The pro Rea Mc Wr Wr Rea Mc	eleness deviation dering deviation ne deviation nich of the following is disadvantage of active replication nisistent replicas nimal bandwidth costs e of the parameters is small nocessing power e effect of a write operation by a process on data item x will always be seen by a successive read operation on x by the same objects. ad-your-writes notonic-writes notonic-writes notonic-Reads ites-follow-reads nich Consistency is used to guarantee that users of a network newsgroup see a posting of a reaction to an article only after the version services ad-your-writes notonic-writes notonic-writes notonic-writes
Nu Sta Oro Tin Wh Co Mi Size Pro The pro Rea Mc Wr Wh hav Rea Mc Mc	eleness deviation dering deviation ne deviation nich of the following is disadvantage of active replication nisistent replicas nimal bandwidth costs e of the parameters is small ocessing power e effect of a write operation by a process on data item x will always be seen by a successive read operation on x by the same ocess. ad-your-writes onotonic-writes onotonic-writes onotonic-reads nich Consistency is used to guarantee that users of a network newsgroup see a posting of a reaction to an article only after the ead-your-writes ad-your-writes onotonic-writes onotonic-writes onotonic-writes onotonic-writes onotonic-writes onotonic-writes onotonic-writes onotonic-writes onotonic-writes
Nu Sta Orc Tin Wh Co Mi Siz Pro Rea Mo Wr Wh hav Rea Mo Mo	eleness deviation dering deviation ne deviation nich of the following is disadvantage of active replication nisistent replicas nimal bandwidth costs e of the parameters is small nocessing power e effect of a write operation by a process on data item x will always be seen by a successive read operation on x by the same objects. nad-your-writes notonic-writes notonic-writes notonic-Reads ites-follow-reads nich Consistency is used to guarantee that users of a network newsgroup see a posting of a reaction to an article only after the version services ad-your-writes notonic-writes notonic-writes notonic-writes
Nu Sta Oro Tin Wh Co Mi Size Pro The pro Rea Mo Wr Wh hav Mo Wr	eleness deviation dering deviation ne deviation nich of the following is disadvantage of active replication nisistent replicas nimal bandwidth costs e of the parameters is small ocessing power e effect of a write operation by a process on data item x will always be seen by a successive read operation on x by the same ocess. ad-your-writes onotonic-writes onotonic-writes onotonic-writes onotonic-Reads rites-follow-reads nich Consistency is used to guarantee that users of a network newsgroup see a posting of a reaction to an article only after the ve seen the original article ad-your-writes onotonic-writes
Nu Sta Orc Tin Wh Co Mi Size Pro The pro Rea Mo Wr Whav Hav Mo Wr The	eleness deviation dering deviation ne deviation nich of the following is disadvantage of active replication nisistent replicas nimal bandwidth costs e of the parameters is small ocessing power e effect of a write operation by a process on data item x will always be seen by a successive read operation on x by the same ocess. ad-your-writes onotonic-writes onotonic-writes onotonic-writes onotonic-Reads rites-follow-reads nich Consistency is used to guarantee that users of a network newsgroup see a posting of a reaction to an article only after the ve seen the original article ad-your-writes onotonic-writes
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Nu Sta Ord Tim Who Color Mile Size Pro Mc	eleness deviation dering deviation ne deviation nich of the following is disadvantage of active replication nisistent replicas nimal bandwidth costs e of the parameters is small decessing power e effect of a write operation by a process on data item x will always be seen by a successive read operation on x by the same ocess. ad-your-writes decorrected by the same ocess. ad-your-writes decorrected by the same ocess. ad-your-writes decorrected by the same ocess. and occorrected by the same occorrecte
Nu Sta Orc Tin Wh Col Mi Size Pro Mc Mc Wr Wh have Rea Mc Wr The effe Rea Mc Mc Wr The Effe Rea Mc	eleness deviation dering deviation ne deviation nich of the following is disadvantage of active replication nisistent replicas nimal bandwidth costs e of the parameters is small nocessing power e effect of a write operation by a process on data item x will always be seen by a successive read operation on x by the same occess. nad-your-writes notonic-writes notonic-writes nich Consistency is used to guarantee that users of a network newsgroup see a posting of a reaction to an article only after the ve seen the original article notonic-writes
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(a)	A crash occurs after sending the completion message and printing the text.
(b)	A crash happens after sending the completion message, but before the text could be printed.
(c)	A crash happens before the server could do anything.
(d)	The text printed, after which a crash occurs before the completion message could be sent.
78	Write the event of following notations: M ~C (~P)
(a)	A crash occurs after sending the completion message and printing the text.
(b)	A crash happens after sending the completion message, but before the text could be printed.
(c)	A crash happens before the server could do anything.
(d)	The text printed, after which a crash occurs before the completion message could be sent.
79	DNS is an example of
(a)	Continuous Consistency
(b)	Causal Consistency
	·
(c)	Sequential Consistency
(d)	Eventual Consistency
80	Staleness deviations relate to the
(a)	first time a replica was updated
(b)	first time a replica was not updated
(c)	last time a replica was updated
(d)	last time a replica was not updated
81	cover the case of dealing with arbitrary failures by process Q, yet these failures are benign: they cannot do any harm.
(a)	Fail-stop failures
(b)	Fail-noisy failures
(c)	fail-silent failures
(d)	Fail-safe failures
00	In which type of failures we assume that communication links are nonfaulty, but that process P cannot distinguish crash failures
82	from omission failures
(a)	Fail-stop failures
(b)	Fail-noisy failures
(c)	fail-silent failures
(d)	Fail-safe failures
83	What are the advantages of file replication?
(a)	Improves availability &performance
(b)	Decreases performance
	Improves consistency
(c)	·
(d)	Improves speed
84	For which file transfer model, cache management is harder due to the variable-length data for different access requests.
(a)	File-level transfer model
(b)	Block-level transfer model
(c)	Byte-level transfer model
(d)	Record-level transfer model
85	Advantage of Data-caching model over the Remote service model
(a)	Reduces network traffic
(b)	Contention for the file servers
(c)	Increased performance
(d)	Contention for Network
86	What is the advantage of caching in remote file access?
(a)	Reduced network traffic by retaining recently accessed disk blocks
(b)	Faster network access
(c)	Copies of data creates backup automatically
(d)	Copies of data resolves inconsistencies
87	In the case of changes to an open file are only locally visible.
(a)	Mutable Files
(b)	Immutable Files
(c)	Atomic Transactions
(d)	Session Semantics
88	Season Schmidts Remote service model for file accessing has
	Reduced contention for the network
(a)	
(b)	Reduced network traffic
(c)	Client's request is performed at the server's node
(d)	No contention for the file servers
89	A cache in which location is the easiest to implement
(a)	Server's main memory
(b)	Client's disk
(c)	Client's main memory
(d)	Server's Disk
90	In a distributed file system, when a file's physical storage location changes
(a)	File name needs to be changed
(b)	File name need not to be changed
(c)	File's host name needs to be changed
(d)	File's local name needs to be changed
91	What are the characteristics of transaction semantics?
(a)	The users of this model are interested in the atomicity property for their transaction
(b)	Suitable for applications that are concerned about coherence of data
(c)	Easy to implement in a single processor system
(d)	Write-back enhances access performance
92	What is NOT the characteristics of session semantics?
	Each client obtains a working copy from the server
(a)	
(b)	When file is closed, the modified file is copied to the file server
(c)	The burden of coordinating file sharing is ignored by the system
(d)	Easy to implement in a single processor system
(-/	

identifies the names of those enterprise-level networks around the world that are linked together via phone, satellite, other communication systems
'
Global Naming Service Directory Services
CODA Files
Andrew File System Which file system implements the most stringent semantics of file sharing
Unix Semantics
Session Semantics
Transaction Semantics
Network Semantics
Which one is NOT a part of the sets of services underlying the Google Search Engine
, , , , , ,
Crawling
Testing
Ranking
Indexing
acts between name servers and their clients.
Name machine
Context node
Naming agents Main agents
Main servers
Which cache location gives Maximum performance gain for file systems
Server's main memory
Client's main memory
Server's Disk
Client's disk
DIT in Domain Name Space stands for
Direct Information Transparency
Directory Implicit Transfer
Direct Information Token
Directory Information Tree
A file appears to the file server as an ordered sequence of records in
Unstructured files
Structured files
Immutable files
Mutable files
Web-based office suite supporting shared editing of documents held on Google servers.
Gmail
Google Talk
Google News
Google Docs
In NFS, all client-server communication is done through
RPC
RMI
Message Oriented Communication
Streams
A system is fault tolerant if it can continue to operate in the
presence of failures
absence of failures
presence of replication
absence of replication